CMPS 101 Summer 2008 Homework Assignment 1

1. (1 Point) p.27: 2.2-2

Consider sorting *n* numbers stored in array *A* by first finding the smallest element of *A* and exchanging it with the element in A[1]. Then find the second smallest element of *A* and exchange it with A[2]. Continue in this manner for the first n-1 elements of *A*. Write pseudo-code for this algorithm, which is known as *selection sort*. What loop invariant does this algorithm maintain? Why does it need to run for only the first n-1 elements, rather than for all *n* elements? Give the best-case and worst-case running times of selection sort in Θ -notation.

2. (1 Point) p.37: 2.3-5

Referring back to the searching problem (see Exercise 2.1-3), observe that if the sequence A is sorted, we can check the midpoint of the sequence against v and eliminate half of the sequence from further consideration. *Binary search* is an algorithm that repeats this procedure, halving the size of the remaining portion of the sequence each time. Write pseudo-code, either iterative or recursive, for binary search. Argue that the worst-case running time of binary search is $\Theta(\lg n)$.

3. (4 Points) p.39: 2-4abcd

Let $A[1 \cdots n]$ be an array of *n* distinct numbers. If i < j and A[i] > A[j], then the pair (i, j) is called an *inversion* of *A*.

- a. (1 Point) List the five inversions of the array (2, 3, 8, 6, 1).
- b. (1 Point) What array with elements from the set {1, 2, 3,, *n*} has the most inversions? How many inversions does it have?
- c. (1 Point) What is the relationship between the running time of insertion sort and the number of inversions in the input array? Justify your answer.
- d. (1 Point) Give an algorithm that determines the number of inversions in any permutation of n elements in $\Theta(n \lg n)$ worst-case time. (Hint: Modify merge sort.)