## **CMPS 101**

## **Final Review Problems**

- 1. Let T be a binary tree, and let N(T) and H(T) denote its number of nodes and height, respectively. Show that  $H(T) \ge |\lg(N(T))|$ . (Hint: this was proved in the solutions to hw6.)
- 2. Let *T* be an almost complete binary tree on *n* nodes, and let *h* be an integer in the range  $0 \le h \le \lfloor \lg(n) \rfloor$ .

Show that the number of nodes in T at height h is exactly  $\left\lfloor \frac{n}{2^h} \right\rfloor - \left\lfloor \frac{n}{2^{h+1}} \right\rfloor$ .

- 3. Trace HeapSort on the following arrays
  - a. (9, 3, 5, 4, 8, 2, 5, 10, 12, 2, 7, 4)
  - b. (5, 3, 7, 1, 10, 12, 19, 24, 5, 7, 2, 6)
  - c. (9, 8, 7, 6, 5, 4, 3, 2, 1)
- 4. Draw the Binary Search Tree resulting from inserting the keys: 5 8 3 4 6 1 9 2 7 (in that order) into an initially empty tree. Write pseudo-code for the following recursive algorithms, and write their output when run on this tree.
  - a. InOrderTreeWalk()
  - b. PreOrderTreeWalk()
  - c. PostOrderTreeWalk()
- 5. Let *T* be a Binary Search Tree and let *x* be a node in *T*. Prove that if *x* has no right child, and if *x* has a successor *y*, then *y* is the lowest ancestor of *x* whose left child is also an ancestor of *x*.
- 6. The predecessor of a node x in a Binary Search Tree is defined to be the node which is printed immediately before x in an InOrderTreeWalk(). Let T be a Binary Search Tree and let x be a node in T. Suppose x has no left child, and x has a predecessor y. State a characterization of the predecessor y similar to the characterization of successor in problem 5. Write an algorithm called TreePredecessor() that returns the predecessor of a non-nil node x, if it exists, and returns nil otherwise.
- 7. State the Binary Search Tree Properties.