**CMPS 201**
**Algorithms and Abstract Data Types**

# Some Common Functions

We present several common functions and estimates which occur frequently in the analysis of algorithms.

**Floors and Ceilings**
Given $x \in \mathbf{R}$, we denote by $\lfloor x \rfloor$ and $\lceil x \rceil$ the *floor of* $x$ and the *ceiling of* $x$, respectively. These are defined to be the unique integers satisfying

$$x-1 < \lfloor x \rfloor \le x \le \lceil x \rceil < x+1$$

Equivalently, if $x \in \mathbf{R}$ and $N \in \mathbf{Z}$ then

(1) $N = \lfloor x \rfloor$ if and only if $N \le x < N+1$, and
(2) $N = \lceil x \rceil$ if and only if $N-1 < x \le N$.

In other words:

(1) $\lfloor x \rfloor$ is the *greatest integer less than or equal to x*, and
(2) $\lceil x \rceil$ is the *least integer greater than or equal to x*.

**Lemma 1:** Let $x \in \mathbf{R}$ and $a, b \in \mathbf{Z}$. Then
(1) $a \le x < b$ if and only if $a \le \lfloor x \rfloor < b$, and
(2) $a < x \le b$ if and only if $a < \lceil x \rceil \le b$.

**Proof of (1):**
(i) $a \le x$ implies $a \le \lfloor x \rfloor$, since among all integers that are less than or equal to $x$, $\lfloor x \rfloor$ is the greatest.
(ii) $x < b$ implies $\lfloor x \rfloor < b$, since $\lfloor x \rfloor \le x$.
(iii) $a \le \lfloor x \rfloor$ implies $a \le x$, since $\lfloor x \rfloor \le x$.
(iv) $\lfloor x \rfloor < b$ implies $x < b$, since $b \le x$ implies $b \le \lfloor x \rfloor$, by (i). ///

**Exercise:** prove part (2).

**Lemma 2:** Let $x \in \mathbf{R}$ and $m \in \mathbf{Z}^+$. Then
(1) $\left\lfloor \frac{\lfloor x \rfloor}{m} \right\rfloor = \left\lfloor \frac{x}{m} \right\rfloor$, and
(2) $\left\lceil \frac{\lceil x \rceil}{m} \right\rceil = \left\lceil \frac{x}{m} \right\rceil$.

**Proof of (1):** Let $N = \lfloor \lfloor x \rfloor / m \rfloor$. Then

$$\begin{aligned} & N \le \frac{\lfloor x \rfloor}{m} < N+1 \\ \Rightarrow\ & mN \le \lfloor x \rfloor < m(N+1) \\ \Rightarrow\ & mN \le x < m(N+1) \qquad \text{(by lemma 1)} \\ \Rightarrow\ & N \le x/m < N+1 \\ \Rightarrow\ & N = \lfloor x/m \rfloor, \end{aligned}$$

and therefore $\lfloor \lfloor x \rfloor / m \rfloor = N = \lfloor x/m \rfloor$. ///

**Exercise:** prove part (2).

<u>**Lemma 3:**</u> Let $a, b, n \in \mathbf{Z}^+$. Then

(1) $\left\lfloor \dfrac{\lfloor n/a \rfloor}{b} \right\rfloor = \left\lfloor \dfrac{n}{ab} \right\rfloor$, and

(2) $\left\lceil \dfrac{\lceil n/a \rceil}{b} \right\rceil = \left\lceil \dfrac{n}{ab} \right\rceil$.

**Proof:** Set $x = n/a$ and $m = b$ in lemma 2. ///

**Exercise**

Let $n \in \mathbf{Z}$. Show that (a) $\left\lfloor \dfrac{n}{2} \right\rfloor + \left\lceil \dfrac{n}{2} \right\rceil = n$, (b) $\left\lceil \dfrac{n}{2} \right\rceil = \left\lfloor \dfrac{n+1}{2} \right\rfloor$, and (c) $\left\lfloor \dfrac{n}{2} \right\rfloor = \left\lceil \dfrac{n-1}{2} \right\rceil$.

**Logarithms**

Let $x, a, b \in \mathbf{R}$ where $x > 0$, $a > 1$, and $b > 1$. Then $\log_a(x)$ denotes the exponent on a which gives *x*. In other words, $\log_a(x)$ is the inverse function of $a^x$, which means $a^{\log_a(x)} = x$ and $\log_a(a^x) = x$. Thus

$$x = a^{\log_a(x)} = \left(b^{\log_b(a)}\right)^{\log_a(x)} = b^{\log_b(a) \cdot \log_a(x)}$$

Taking $\log_b$ of both sides of this equation yields

(*) $$\log_b(x) = \log_b(a) \cdot \log_a(x),$$

which says in particular $\log_b(x) = \text{constant} \cdot \log_a(x)$, i.e. any two log functions differ by a constant multiple. It follows that $\log_b(n) = \Theta(\log_a(n))$, so speaking in terms of asymptotic growth rates, there is really only one log function. Equation (*) implies

$$\log_a(x) = \frac{\log_b(x)}{\log_b(a)}$$

which shows how to convert from one log function to another. In particular $\lg(x) = \frac{\ln(x)}{\ln(2)}$. Here we use the standard notation $\lg(\,) = \log_2(\,)$, and $\ln(\,) = \log_e(\,)$, where $e = 2.71828..$ . Equation (*) also implies $a^{\log_b(x)} = a^{\log_a(x)\cdot\log_b(a)} = \left(a^{\log_a(x)}\right)^{\log_b(a)} = x^{\log_b(a)}$, which gives us the useful formula

$$a^{\log_b(x)} = x^{\log_b(a)}.$$

**Stirling's Formula**

Let $n \in \mathbf{Z}^+$. Then $n! = \sqrt{2\pi n} \cdot \left(\frac{n}{e}\right)^n \cdot \left(1 + \Theta\left(\frac{1}{n}\right)\right)$.

Stirling's formula gives a simple way to determine asymptotic (upper, lower, and tight) bounds on functions involving $n!$. An elementary proof can be found at

http://www.sosmath.com/calculus/sequence/stirling/stirling.html

**Corollary:**

(1) $n! = o(n^n)$

(2) $n! = \omega(b^n)$ for any $b > 0$

(3) $\log(n!) = \Theta(n\log(n))$

**Proof of (1):**

$$\frac{n!}{n^n} = \frac{\sqrt{2\pi n} \cdot \left(\frac{n}{e}\right)^n \cdot \left(1 + \Theta\left(\frac{1}{n}\right)\right)}{n^n} = \frac{\sqrt{2\pi n} \cdot \left(1 + \Theta\left(\frac{1}{n}\right)\right)}{e^n} \to 0 \text{ as } n \to \infty$$, showing that $n! = o(n^n)$. ///

**Proof of (3):** Taking log (any base) of both sides of Stirling's formula, we get

$$\log(n!) = \log\sqrt{2\pi n} + \log\left(\frac{n}{e}\right)^n + \log\left(1 + \Theta\left(\frac{1}{n}\right)\right)$$

$$= \frac{1}{2}\log(2\pi) + \frac{1}{2}\log(n) + n\log(n) - n\log(e) + \log\left(1 + \Theta\left(\frac{1}{n}\right)\right).$$

Therefore

$$\frac{\log(n!)}{n\log(n)} = 1 + \left(\text{stuff that} \to 0 \text{ as } n \to \infty\right),$$

hence $\lim_{n\to\infty}\left(\frac{\log(n!)}{n\log(n)}\right) = 1$, proving that $\log(n!) = \Theta(n\log(n))$. ///

**Exercise:** Prove part (2) of the corollary.

**Exercise:** Prove that $\binom{2n}{n} = \Theta\left(\frac{4^n}{\sqrt{n}}\right)$, where $\binom{m}{k}$ denotes the binomial coefficient $\binom{m}{k} = \frac{m!}{k!(m-k)!}$, for $0 \le k \le m$.