

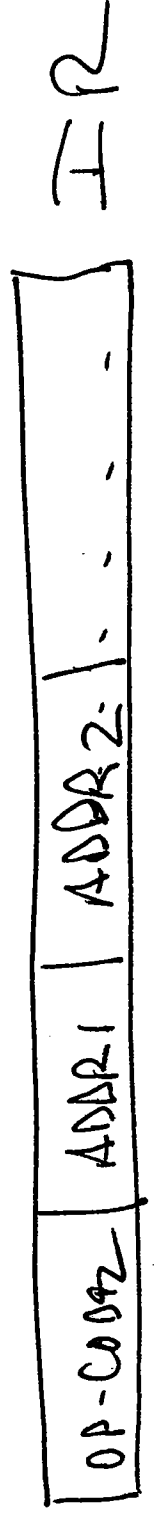
CNAS 10 3-6-08



The Control Unit

- 1.) Fetch next instruction from memory
- 2.) Decode the instruction
- 3.) Execute the instruction.

Instruction Register



$$\therefore \text{Max \# of operations} = 2^8 = 256$$

$$\therefore \text{Max size of memory unit} = 2^{24} \text{ Bytes} \\ = 16 \text{ MB}$$

$$\left(\text{why } 2^{24} = \underbrace{2 \cdot 2}_{16}^{20} = 16 \text{ MB} \right)$$

\therefore # bytes needed to store 1 instruction

$$= 10$$

4



= 32 Bits

4 14 14

Ex. IR

- max # of OAS = $16 = 2^4$
- max size of memory = 2^{14} Bytes
= 16 KB
- # Bytes needed to store one instruction = 4

Instruction Set : The set of all

operations that can be performed by
a processor.

CISC : Complex instruction set architectures
200 - 400 ops.

RISC : Reduced instruction set architectures
30 - 50 ops

EX

OP-CODE	ADDR-1	ADDR-2	IR
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 $8 \times 8 \times 8 = 24 \text{ bits}$

SUPPOSE WE WISH TO MOVE CONTENTS
 OF CELL 27 TO CELL 53, AND
 SUPPOSE THE OP-CODE FOR THIS DATA
 TRANSFER IS 15

NOTE: $15 = 1111$, $27 = 11011$, $53 = 110101$

OP-CODE ADDR-1 ADDR-2

0	0	0	0	1	1	1	1	0	0	0	1	1	0	1	1	0	0	1	1	0	1	0	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

= R

STANDARDS: MOVE 27 53
" " "

00001111 00011011 00110101

NOTATIONAL CONVENTIONS

- X, Y, Z, ... STAND FOR CELL ADDRESSES.
- a, b, c, ... STAND FOR CELL CONTENTS.
- ASSUMED ALU CONTAINS JUST ONE REGISTER: R

□

• WRITE CON(X) TO MEMORY CONTAINS OF CELL X'

FOL CLASSES OF MACHINE LANGUAGE INSTRUCTIONS

1.) DATA TRANSFER.

- MOVE X Y : CON(Y) ← CON(X)
- LOAD X : CON(R) ← CON(X)
- STORE X : CON(X) ← CON(R)

2.) ARITHMETIC

- ADD₁ X : CON(R) ← CON(X) + CON(R)
- ADD₂ X Y : CON(Y) ← CON(X) + CON(Y)
- ADD₃ X Y Z : CON(Z) ← CON(X) + CON(Y)

⋮

3.) COMPARISON

SET VALUES OF SPECIAL BITS CALLED

COMPARISON CODES : GT, EQ, LT

LD



In ALU:

o COMPARE X Y : COMPARES CON(X) TO
 CON(Y) AND SET
 CONDITION CODES AS

Follows :

IF
 $CON(X) > CON(Y)$
 $CON(X) = CON(Y)$
 $CON(X) < CON(Y)$

GT	EQ	LT
1	0	0
0	1	0
0	0	1

4.) BRANCHING (ALTER NORMAL SEQ. FLOW OF PROGRAM EXECUTION)

• JUMP X : TAKE NEXT INSTRUCTION

FROM CELL X,

• JUMPGT X : IF $GT = 1$, TAKE NEXT

INSTR. FROM CELL X, OTHERWISE
FOLLOW SEQUENTIAL FLOW
OF EXECUTION.

• JUMPED X

• JUMPLT X

• JUMPLE X

• JUMPGRE X

• JUMPNEQ X

o HALT

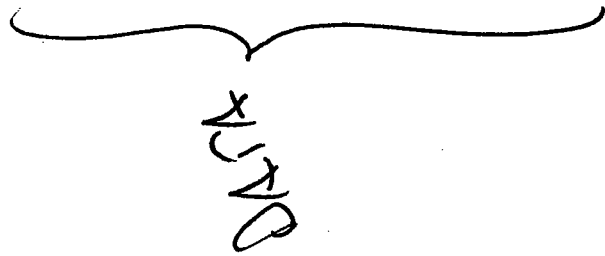
:Quit

Fetch-Decode-Execute.

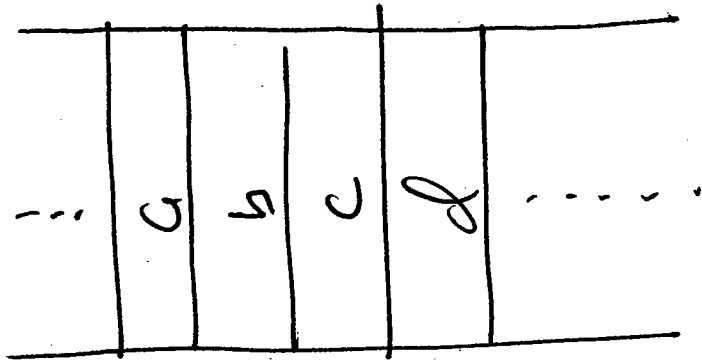
In following example Assume :

address

200
201
202
203



memory



Also Assume
PROGRAM INSTRUCTIONS
BEGIN AT CELL 100

100
101

